

UTILITY
PATENT APPLICATION
TRANSMITTAL

(Only for new nonprovisional applications under 37 CFR
1.53(b))

12-26-00
A
Attorney Docket No. 1400.4100290 Total Pages 42
First Inventor or Application Identifier McCormick, et al.
Title BUFFERING SYSTEM FOR USE IN A
COMMUNICATION SWITCH THAT INCLUDES A
MULTIPROCESSOR CONTROL BLOCK AND
METHOD THEREFORE
Express Mail Label No. EL454269848US

JC853 U.S. PTO
09/17/00
12/21/00

APPLICATION ELEMENTS See MPEP chapter 600 concerning utility patent application contents.	ADDRESS TO: Assistant Commissioner for Patents Box Patent Application Washington, DC 20231
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1. ☒ Fee Transmittal Form
(Submit an original, and a duplicate for fee processing)

2. ☒ Specification *Total Pages 28*
(preferred arrangement set forth below)
- Descriptive title of the Invention
- Cross References to Related Applications
- Statement Regarding Fed sponsored R & D
- Reference to Microfiche Appendix
- Background of the Invention
- Brief Summary of the Invention
- Brief Description of the Drawings (if filed)
- Detailed Description
- Claim(s)
- Abstract of the Disclosure

3. ☒ Drawings (35 USC 113) *Total Sheets 7*

4. Oath or Declaration *Total Pages 2*
a. ☒ Newly executed (original or copy)
b. ☐ Copy from a prior application
(37 CFR 1.63(d))
(for continuation/divisional with Box 16 completed)

i. ☐ **DELETION OF INVENTOR(S)**
Signed statement attached deleting
inventor(s) named in the prior application,
see 37 CFR 1.63(d)(2) and 1.33(b).

5. ☐ Microfiche Computer Program (Appendix)
6. ☐ Nucleotide and/or Amino Acid Sequence Submission
(if applicable, all necessary)

a. ☐ Computer Readable Copy
b. ☐ Paper Copy (identical to computer copy)
c. ☐ Statement verifying identity of above copies

ACCOMPANYING APPLICATION PARTS

7. ☒ Assignment Papers (cover sheet & document(s))
8. ☐ 37 CFR 3.73(b) Statement ☐ Power of
(when there is an assignee) Attorney
9. ☐ English Translation Document (if applicable)
10. ☐ Information Disclosure ☐ Copies of
Statement (IDS)/PTO-1449 IDS Citations
11. ☐ Preliminary Amendment
12. ☒ Return Receipt Postcard (MPEP 503)
(Should be specifically itemized)
13. ☐ Small Entity ☐ Statement filed in Prior
Statement(s) Application, Status still proper and
desired.
14. ☐ Certified Copy of Priority Document(s)
(if foreign priority is claimed)
15. ☐ Other

16. If a CONTINUING APPLICATION, check appropriate box and supply the requisite information:

☐ Continuation ☐ Divisional ☐ Continuation-in-part (CIP) of prior application No:

Prior Application Information: Examiner

Group / Art Unit:

For CONTINUATION or DIVISIONAL APPS only: The entire disclosure of the prior application, from which an oath or declaration is supplied under Box 4B, is considered a part of the disclosure of the accompanying continuation or divisional application and is hereby incorporated by reference. The incorporation can only be relied upon when a portion has been inadvertently omitted from the submitted application parts.

17. CORRESPONDENCE ADDRESS

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Signature	Ross D. Snyder	Date	08 December 2000

FEE TRANSMITTAL

Note: Effective October 1, 1997.
Patent fees are subject to annual revision.

TOTAL AMOUNT OF PAYMENT (\$) 804.00

METHOD OF PAYMENT (check one)

1. ☒ The Commissioner is hereby authorized to charge indicated fees and credit any over payments to:

Deposit Account Number	50-1566
Deposit Account Name	Ross D. Snyder & Associates, Inc.

☒ Charge Any Additional Fee Required Under 37 CFR 1.16 and 1.17

☐ Charge the Issue Fee Set in 37 CFR 1.18 at the mailing of the Notice of Allowance

2. ☒ Payment Enclosed:

☒ Check ☐ Money Order ☐ Other

FEE CALCULATION

1. FILING FEE

Large Entity Fee Code (\$)	Small Entity Fee Code (\$)	Fee Description	Fee Paid
101 710	201 355	Utility filing fee	710.00
106 320	206 160	Design filing fee	
107 4900	207 245	Plant filing fee	
108 710	208 355	Reissue filing fee	
114 150	214 75	Provisional filing fee	

SUBTOTAL (1) (\$) 710.00

2. CLAIMS

Claims	Extra	Fee from below	Fee Paid
Total 23	(-20 =) 3	18.00	54.00
Indep. 3	(-3 =) 0		
Multiple Dep.			

Large Entity Fee Code (\$)	Small Entity Fee Code (\$)	Fee Description
103 18	203 9	Claims in excess of 20
102 80	202 40	Independent claims in excess of 3
104 270	204 135	Multiple dependent claim
109 80	209 40	Reissue independent claims over original patent
110 18	210 9	Reissue claims in excess of 20 and over original patent

SUBTOTAL (2) (\$) 54.00

Complete if Known

Application Number	
Filing Date	08-10-2000
First Named Inventor	McCormick, et al.
Group Art Unit	
Examiner Name	
Attorney Docket Number	1400.4100290

FEE CALCULATION (continued)

3. ADDITIONAL FEES

Large Entity Fee Code (\$)	Small Entity Fee Code (\$)	Fee Description	Fee Paid
105 130	205 65	Surcharge - late filing fee or oath	
127 50	227 25	Surcharge - late provisional filing fee or cover sheet	
139 130	139 130	Non-English specification	
147 2,520	147 2,520	For filing a request for reexamination	
112 920*	112 920*	Requesting publication of SIR prior to Examiner action	
113 1,840*	113 1,840*	Requesting publication of SIR after Examiner action	
115 110	215 55	Extension for reply within first month	
116 390	216 195	Extension for reply within second month	
117 890	217 445	Extension for reply within third month	
118 1,390	218 645	Extension for reply within fourth month	
128 1,890	228 945	Extension for reply within fifth month	
119 310	219 155	Notice of Appeal	
120 310	220 1550	Filing a brief in support of an appeal	
121 270	221 135	Request for oral hearing	
138 1,510	138 1,510	Petition to institute a public use proceeding	
140 110	240 55	Petition to revive - unavoidable	
141 1,240	241 620	Petition to revive - unintentional	
142 1,240	242 620	Utility issue fee (or reissue)	
143 440	243 220	Design issue fee	
144 600	244 300	Plant issue fee	
122 130	122 130	Petitions to the Commissioner	
123 50	123 50	Petitions related to provisional applications	
126 240	126 240	Submission of Information Disclosure Stmt	
581 40	581 40	Recording each patent assignment per property (times number of properties)	40.00
146 710	246 355	Filing a submission after final rejection (37 CFR 1.129(a))	
149 710	249 355	For each additional invention to be examined (37 CFR 1.129(b))	
Other fee (specify)			
Other fee (specify)			

* Reduced by Basic Filing Fee Paid

SUBTOTAL (3) (\$) 40.00

SUBMITTED BY ROSS D. SNYDER & ASSOCIATES, INC.

Complete (if applicable)

Typed or Printed Name

Ross D. Snyder, Reg. No. 37,730

Signature

Ross D. Snyder

Date

08 Dec 00

Deposit Account User ID

PATENT APPLICATION

IN THE UNITED STATES PATENT AND TRADEMARK OFFICE

Applicant: McCormick, et al.

Examiner:

Serial No:

Art Group:

Filing Date: 08-10-00

Docket No: 1400.4100290

Title: BUFFERING SYSTEM FOR USE IN A COMMUNICATION SWITCH THAT INCLUDES A MULTIPROCESSOR CONTROL BLOCK AND METHOD THEREFORE

To the Honorable Commissioner
of Patents and Trademarks
Washington, D.C. 20231

CERTIFICATE OF EXPRESS MAILING

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Name of Depositor: **Terri Alloway**
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Date of Deposit: 12-20-2000

Signature: Terri Alloway

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Statement of Small Entity Status;



Information Disclosure Statement;



Return upon receipt post card;

PATENT APPLICATION
DOCKET NO. 1400.4100290

In the United States Patent and Trademark Office

FILING OF A UNITED STATES PATENT APPLICATION

Title:

BUFFERING SYSTEM FOR USE IN A COMMUNICATION SWITCH THAT
INCLUDES A MULTIPROCESSOR CONTROL BLOCK AND METHOD
THEREFORE

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Date of Deposit: 12-20-2000

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Name of Depositor: TERRI ALLOWAY
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Signature: Terri Alloway

**BUFFERING SYSTEM FOR USE IN A COMMUNICATION SWITCH THAT
INCLUDES A MULTIPROCESSOR CONTROL BLOCK AND METHOD
THEREFORE**

5

Related Applications

This application is claims priority to a provisional application 60/224,441 filed August 10, 2000 having the same title as the present application. The present application is related to a co-pending application entitled "MULTIPROCESSOR CONTROL BLOCK FOR USE IN A COMMUNICATION SWITCH AND METHOD

10 THEREFORE", which has an attorney docket number of 1400.4100220 and which was filed on the same date as the present application.

Background of the Invention

Switches commonly used in communication networks often include a plurality of line cards that are intercoupled by a switching fabric. Each of the plurality of line cards receives ingress data and transmits egress data. The switching fabric allows ingress data received by one card to be directed to an egress connection on one or more of the other line cards included within the communication switch.

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Various processing operations are typically performed for the various connections supported by the communication switch. Such connections can include both switched and permanent connections. The processing operations include routing that must occur within the switch such that ingress data received by one line card is directed to the appropriate egress line card. In prior art systems, the processing operations including routing were typically performed by individual processors included in each of the line cards of the communication switch. Figure 1 illustrates a prior art switch that includes a switching fabric and a plurality of line cards. As is illustrated, each of the line cards includes a processor that is responsible for performing the various call processing and routing functions for connections that are directed through that particular line card.

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Such distributed call processing provides some disadvantages that become more apparent as data communication speeds continue to increase. In some cases, specific line cards within the switch will be much more active than other line cards. Some of the more active line cards can become overwhelmed by the amount of call traffic they are required to service. Because of the distributed nature of the call processing amongst the various line cards, a processor on a line card that is not being fully utilized is unable to assist the overburdened processor on another line card. As such, the overall call processing bandwidth available within the switch is not fully utilized.

Another disadvantage of distributed call processing is that none of the processors in the individual line cards has a global perspective on the operation of the switch. As such, the adaptations that an individual processor is capable of performing in order to increase the efficiency of the operation of the switch are limited.

Therefore, a need exists for a communication switch that includes call processing and routing functionality that does not suffer from the disadvantages presented by distributed processing of prior art systems.

Brief Description of the Drawings

Figure 1 illustrates a block diagram of a prior art switch;

Figure 2 illustrates a block diagram of a communication switch that includes a multiprocessor control block in accordance with a particular embodiment of the present invention;

Figure 3 illustrates a block diagram of one embodiment of the multiprocessor control block of Figure 2;

Figure 4 provides a graphical representation of a protocol stack associated with processing within a communication switch in accordance with a particular embodiment of the present invention;

Figure 5 illustrates a block diagram of an alternate embodiment of the
5 multiprocessor control block of Figure 2;

Figure 6 illustrates a block diagram of the various components included in a communication switch that includes a multiprocessor control block in accordance with a particular embodiment of the present invention; and

Figure 7 illustrates a block diagram showing the various queuing structures
10 included in the various processing entities within a communication switch that includes a multiprocessor control block in accordance with a particular embodiment of the present invention.

Detailed Description

Generally, the present invention provides a buffering system in a communication
15 switch and method therefore that utilizes notification of congested points within the switch to perform intelligent routing decisions such that the congestion is avoided when possible. In one embodiment, this is accomplished by detecting congestion in a transmit queue corresponding to a line card or a link layer processor of the communication switch. Upon detection of such congestion, an indication of the congestion is provided to a
20 central control block that performs call processing and routing for the communication switch. When the central control block performs subsequent routing operations, the central control block considers the congestion indications it has thus far received such that the congested points within the switch are avoided if possible.

By including a central control block within the switch that oversees all of the routing for the switch, the notification of various congested points within the switch can be used to find alternate routes. This was not possible in prior art systems where distributed call processing and routing processors within the various line cards did not possess the global point of view of the central control processor described herein. The present invention also provides techniques for isolating congestion within the switch to predetermined congestion points, where when congestion occurs at these predetermined points, congestion indications are generated and such information is disseminated to the various processors supporting the switch such that intelligent decisions can be made by these processors.

The invention can be better understood with reference to Figures 2-7. Figure 2 illustrates a block diagram of a communication switch 100 that includes a multiprocessor control block 110, a switching fabric 120, and a plurality of line cards 132-138. As is apparent to one of ordinary skill in the art, the number of line cards included within the switch 100 may vary from one implementation to the next. The switch 100 may be a switch used in an ATM network, or some other communication network that supports data communication using other protocols. For example, the switch 100 may support internet protocol (IP) packet traffic, where such packetized traffic may utilize Packet over Sonet (POS), multiprotocol label switching (MPLS), or packet over ATM communication techniques.

Rather than including individual processors within each of the line cards 132-138 to perform call processing and routing functionality as was the case in prior art systems, a unified multiprocessor control block 110 performs all of these functions for all of the line cards 132-138 included within the switch 100. Although this may reduce some of the benefits in terms of modularity and ease of expansion that existed in prior art distributed processing switches, the benefits obtained through the centralized processing greatly outweigh those forfeited.

Each of the line cards 132-138 includes ingress and egress connections. Ingress data traffic is received over the ingress connections and forwarded across the switching fabric 120 to an egress line card where it is sent out over an egress connection as egress data traffic. In some cases, ingress data traffic received by a line card will include call setup messages, or other control data that relates to call processing or routing functionality. Such traffic is forwarded to the multiprocessor control block 110 for further processing. The result of such processing may produce egress data units such as acknowledgement messages, which may be cells or packets that are to be forwarded through an egress line card to another switch of the communications network within which the switch 100 is included.

Figure 3 provides a block diagram of a particular embodiment of a multiprocessor control block 110 that may be included in the switch 100 of Figure 2. The multiprocessor control block 110 includes a resource and routing processor 220, a plurality of intermediate processors 230-234, and a link layer processor 240. In other embodiments, a different array of processors may be included in the multiprocessor control block 110, where a division of functionality amongst the various processors is performed such that efficiency of operation of the multiprocessor control block 110 is optimized. The various functions performed by the processors included in the multiprocessor control block 110 may include those associated with a protocol stack 290 illustrated in Figure 4.

The protocol stack 290 is an example protocol stack typically associated with a data communication system. The physical layer 296 corresponds to the lowest section of the protocol stack, and the physical layer may include operations relating to the specific protocol over which the data is being transmitted within the switch. Such processing can include support for protocols such as ATM, where the ingress and egress data units processed by the multiprocessor control block 110 may include ATM adaptation layer 5 (AAL5) packets made up of multiple ATM cells. In other embodiments, the physical layer 296 may support POS, DWDM, etc.

The functionality included in the physical layer processing 296 may be wholly performed within the link layer processor 240 included within the multiprocessor control block 110 of Figure 3. In many cases, the amount of processing required to support the physical layer 296 is limited such that a single link layer processor 240 can provide all of the support for the data traffic requiring such processing within the switch. Conventional terminology would describe the processing operations performed by the link layer processor 240 as including all of the layer 2 functionality associated with the protocol stack 290.

The other portions of the protocol stack 290 associated with the multiprocessor control block include those corresponding to the signaling link layer 295, signaling protocol 294, call processing 293, and connection resource allocation and routing 292. The signaling link layer portion 295 includes functionality such as verification that the various links required by the system are up. The signaling protocol portion 294 includes support for control traffic required for setting up calls, tearing down calls, etc. Various protocols may be supported by this portion, including MPLS, narrow band ISDN, broadband ISDN, etc. The call processing portion 293 includes the functionality associated with the operations required to support the processing of calls or connections that utilize the switch within their data path. The connection resource allocation and routing portion 292 includes the functionality associated with selecting egress routes for connections that utilize the switch, and allocating the available data bandwidth and connection resources within the switch amongst the multiple calls supported.

In addition to performing all of the processing operations corresponding to layer 2, the link layer processor 240 may also perform some of the processing operations associated with layer 3 processing. Such layer 3 processing may include the distribution of the responsibility for maintaining the call resources (e.g. context information, etc.) for individual calls to the various intermediate processors that are present within the system. This will be described in additional detail below. Similarly, the link layer processor 240 may also perform functions typically associated with layer 2 processing such as the

distribution of global call references or signaling messages that apply to all of the calls that are active. This is also described in additional detail below.

Each of the intermediate processors 230-234 performs call-processing operations for a corresponding portion of the connections supported by the switch. Call processing
5 includes handling subscriber features for various calls, wherein a particular example, such subscriber features may include caller ID, private line, or other features that require the switch to perform certain support tasks. The functionality associated with such call processing is easily distributed amongst multiple processors such that efficiencies can be obtained through parallel processing of multiple requests. As is apparent to one of
10 ordinary skill in the art, if the system does not require parallel processing resources for efficient calling processing, a single intermediate processor or a fewer number of intermediate processors may be included within the multiprocessor control block 110. In other cases, the number of intermediate processors may be increased to supplement the parallel processing capabilities of the multiprocessor control block 110.

15 The resource and routing processor 220 performs the functions associated with resource distribution and routing. By having a global perspective on the available routes and resources within the entire switch, the resource and routing processor is able to perform its duties in a manner that may provide better alternate routes or better utilization of resources that was not possible in prior art systems. The various resources that are
20 allocated by the resource and routing processor 220 may include the channels available between the switches, where in an ATM network, such channels may include virtual channel connections (VCCs) and virtual path connections (VPCs). In an MPLS system, resources may include the distribution of the MPLS labels associated with the label switched paths existing within the network. The bandwidth allocation that may be
25 performed by the resource and routing processor 220 may include performing connection admission control (CAC) operations. The bandwidth available for allocation may be divided up based on the quality of service (QOS), traffic rates, etc. associated with the various lengths within the switch.

Thus, the various portions of the protocol stack 290 shown in Figure 4 are divided amongst the resource and routing processor 220, the plurality of intermediate processors 230, and the link layer processor 240. Such a division allows for a common processor to support certain functionality within the switch whereas when distribution of the processing is possible to multiple processors, the advantages of parallel processing can be realized.

When the link layer processor 240 receives ingress data units that must be processed, such as ingress data units corresponding to a call setup message, the link layer processor 240 forwards those ingress data units to a selected one of the intermediate processors 230-234 for additional processing. As stated above, the intermediate processors 230-234 all perform similar processing operations, where multiple intermediate processors allow for parallel processing to occur. Specifically, the intermediate processors perform functions such as those included in the signaling layer of the protocol stack described with respect to Figure 4 as well as call processing operations.

The link layer processor 240 determines to which of the intermediate processors 230-234 the ingress data units corresponding to a call setup message are forwarded based on some type of prioritization scheme. In some embodiments, the prioritization scheme may be as simple as a round robin scheme such that the link layer processor 240 moves from one intermediate processor to the next as ingress data units corresponding to various messages are received. In other embodiments, the prioritization scheme may take into account the loading on each of the intermediate processors when assigning the ingress data units to one of the intermediate processors for processing. Thus, if one intermediate processor is overloaded, rather than forwarding new data to be processed to that processor, a less loaded intermediate processor is selected instead. The loading on the various intermediate processors may be determined based on the current loading on the ingress queues that receive data units from the link layer processor 240. Such ingress queues within the intermediate processors 230-234 are described in additional detail with respect to Figure 7 below.

At various points within the multiprocessor control block 110, a sequence number that corresponds to a particular call may be assigned to each of the ingress data units for that call. As such, when the intermediate processor completes its processing operations on the ingress data unit and forwards it to the resource and routing processor 220 for subsequent processing, the sequence number can be used to ensure that the ordering of the various processing operations is maintained or at least understood. Thus, the resource and routing processor 220 may utilize the sequence numbers to identify a particular point in a processing time line that is associated with each call. If resources corresponding to a particular call are subsequently allocated to a different call and a command is later received by the resource and routing processor 220 corresponding to the previously existing call, that command will include the sequence number associated with the previous call. As such, the resource and routing processor 220 can choose to discard or disregard the command, as the call is no longer active. However, the resource and routing processor can also identify similar commands for the current call that is now utilizing those resources based on a match between the current sequence number associated with the resources addressed in the command, as that sequence number also corresponds to the current call that has been established.

In a specific example, assume that a first command is received that establishes a call using a virtual connection identifier (VCI) 20. A short time later, a decision is made to reallocate that VCI to a new call. In order to reassign the VCI, the old call is released and the new call is created. If the two commands associated with releasing the old call and creating the new call are sent in quick succession to the switch, these commands may take different parallel paths through different intermediate processors. These parallel paths may have different latencies such that the call setup message associated with the VCI 20 may be received by the resource and routing processor prior to the call release message associated with the old call. When the resource and routing processor 220 services the call setup message such that the VCI 20 is now assigned to the new call, it will assign that call a new sequence number. Thus, when the call release command associated with the sequence number of the old call is received, the resource and routing

processor 220 will recognize that this command was intended to release the old call and not the call that has just been set up. As such, the newly set up call will not be released.

In order to ensure that certain processing operations that must be performed by the same intermediate processor for a specific call are performed by the same
5 intermediate processor, the link layer processor 240 may maintain some type of indication as to where certain ingress data units have been forwarded in the past. Thus, if a first message corresponding to a first call is directed to the intermediate processor 230, and a second message corresponding to the same first call is later received that requires processing by the intermediate processor 230, the link layer processor 240 will have
10 maintained a record as to which of the intermediate processors processed the first message such that appropriate forwarding of the second message can occur. The context for specific calls maintained within one of the intermediate processors 230-234. As such, when messages corresponding to the call context are received that may modify or utilize the call context, they must be forwarded to the appropriate intermediate processor for
15 processing.

Thus, for messages corresponding to new calls, or messages that are processed independently of a specific intermediate processor associated with a specific call, the link layer processor 240 may forward such messages based on the prioritization scheme. In other cases where dependencies between a received message and a previously processed
20 message exists, the link layer processor 240 may forward these ingress data packets based on history data that it stores relating to ingress data packets that it has forwarded in the past.

In some instances, global call reference messages directed to all calls will be received by the link layer processor 240. Such global call references or universal
25 signaling messages are recognized by the link layer processor 240 and distributed to all of the intermediate processors 230-234. In some instances, such global call references or signaling messages require a response. When this is the case, the link layer processor

240 may collect responses to the global call reference from each intermediate processor and compile these responses to produce a unified response.

An example of a global call reference or signaling message that applies to all calls is a "clear all calls" message. Such call clearing requires acknowledgement. As such,
5 when the link layer processor 240 receives such a global call reference, it forwards it to all of the intermediate processors 230-234. The link layer processor 240 then collects all of the acknowledgements corresponding to the call clearing performed within the intermediate processors 230-234 and sends an accumulated acknowledgement to the entity which issued the clear all calls message.

10 In order to facilitate interaction between the various intermediate processors 230-234 and the line cards included in the switch, a message processor 250 may be included in the multiprocessor control block 110. The message processor 250, which is operably coupled to the plurality of intermediate processors 230-234 and the plurality of line cards included in the switch, supports messaging between the plurality of intermediate
15 processors 230-234 and one or more of the line cards. The message processor 250 may act as a queuing point for messages between the intermediate processors 230-234 and the line cards, where multiple messages are bundled together for distribution to the various line cards in order to improve the efficiency with which such messages are communicated.

20 The multiprocessor control block 110 may also include a management block 210 that is operably coupled to the other processors included within the multiprocessor control block 110. The management block 210 receives management requests 202 that it processes to produce configuration commands issued to the various processors included within the multiprocessor control block 110. Management block 210 may also provide
25 configuration information to line cards included within the switch. Management block 210 may be used to retrieve statistics or current configuration settings and status information from the various line cards and processors included within the switch.

Furthermore, the management block 210 may store a database that can be used for auto-recovery by the switch in the case of a power outage or similar failure.

The specific multiprocessor control block 110 illustrated in Figure 3 only includes a single link layer processor 240 that supports all of the line cards included within the switch. In another embodiment illustrated in Figure 5, a single resource and routing processor 320 supports two sets of distributed processors associated with different sets of line cards 342, 362 included within the switch. Thus, the first set of line cards within the switch 342 is supported by a link layer processor 340. The link layer processor 340 is coupled to a first plurality of intermediate processors 330-334 that perform the call processing and signaling layer processing for the calls associated with the first set of line cards 342.

Similarly, a second set of line cards 362 is supported by a second link layer processor 360. The link layer processor 360 forwards ingress data units received via the second set of line cards 362 to an appropriate one or more of the intermediate processors 352-356.

Although a greater level of processing distribution exists in the multiprocessor control block illustrated in Figure 5 than that illustrated in Figure 3, a single resource and routing processor 320 is still used in the multiprocessor control block 110 of Figure 5 to oversee all of the resource allocation and routing for the entire switch and all of its corresponding line cards. As such, the benefits of a centralized resource and routing processor are realized along with possibly some additional benefit in the distribution of other portions of the processing tasks associated with the maintenance of the calls within the switch.

In many instances, congestion within a data communication switch can result in the undesired loss of certain messages that may be crucial to the proper operation of the switch. As such, reduction or elimination of congestion within the switch is highly

desirable. By positioning queuing structures at appropriate points within the switch and ensuring that congestion is isolated to certain selected queuing points, overall congestion within the switch can be greatly reduced, thus improving the efficiency with which the switch is able to operate.

5 Figure 6 illustrates a general architectural view of a communication switch that includes a multiprocessor control block. The multiprocessor control block includes the resource and routing processor 610, the plurality of intermediate processors 622-626, the link layer processor 630, and may also include the message processor 640. The message processor 640 allows the plurality of intermediate processors 622-626 to interact with the
10 line cards 652-656 included in the switch. The data flow amongst these various processors and the line cards presents numerous points at which queuing and monitoring of such queuing can be used to reduce congestion and improve throughput within the switch. In order to illustrate the various queuing structures that can be included in the processors and line cards, a simplified block diagram is provided in Figure 7.

15 Figure 7 reduces the number of entities included within the switch such that only a single intermediate processor 622 and a single line card 652 are shown. However, the various queue structures shown for these blocks are preferably also included in additional intermediate processors and line cards that may included within the switch.

As stated above, the queuing structures included in the various processors and
20 other parts of the system are intended to allow the entities to interact with each other in a manner that makes maximum use of the resources available. However, certain entities within the system may be overloaded to the point where there is too much data flowing through the system for the data to be dealt with appropriately. As such, queue backups and congestion can occur. By configuring the system such that congestion is
25 concentrated at specific queue structures, those queue structures can be designed to include special control circuitry that deals with the congestion in an intelligent manner. By isolating the congested queuing points and understanding how those points become

congested, the appropriate measures that may be taken to help alleviate the congestion are better understood. Furthermore, by only putting the circuitry adapted to deal with congestion at those points where congestion is expected to concentrate, the overhead associated with such additional circuitry is minimized.

5 Referring to Figure 7, the link layer processor 630 receives ingress data units at a queuing point 702. These ingress data units may be received from a neighboring node within the network. Because of the link protocol, only a fixed number of ingress data units will be received from the neighboring node until an acknowledgement is sent to the neighboring node that indicates that the ingress data units already sent have been properly
10 processed. Thus, the neighboring node will not send over any more ingress data units until those already sent have been processed. As such, this queuing point is a “windowed” queuing point such that it cannot become congested due to the feedback path that exists.

From the queuing point 702, the ingress data units are routed to a particular
15 intermediate processor of the plurality of intermediate processors included in the multiprocessor control block. The example shown in Figure 7, ingress data units are forwarded from the queuing point 702 to a queuing point 712 included in the intermediate processor 622. The queuing point 712 may have a corresponding threshold level that is used to monitor the contents of the queue 712. If the threshold level is exceeded, an
20 indication is relayed back to the link layer processor 630 and the windowing available at the queuing node 702 is used to slow down the inflow of data units such that the flow to the queuing point 712 is limited until the threshold is no longer exceeded. Thus, the windowing available at the queuing point 702 can be used to ensure that congestion does not build up at the queuing node 712 in the intermediate processor 622. In some
25 embodiments, the windowing function at the queuing node 702 simply discards incoming data units when congestion exists within the intermediate processor 622. In other embodiments, a buffer of limited size may be present at the queuing point 702 such that

when the limits of this limited buffer are exceeded, discarding of the incoming data units occurs until the congestion is alleviated.

After retrieval from the queue 712 and processing by the intermediate processor 622, data units are forwarded to the resource and routing processor 610, where they are stored in a queue 722. Typically, the resource and routing processor 610 has adequate bandwidth to act quickly on the data that it receives at the queuing point 722. As such, congestion is not typically a problem within the resource and routing processor 610. However, in the case where conditions cause congestion, a similar threshold determination within the queue 722 can trigger a processing stoppage in one or more of the plurality of intermediate processors 622. This may cause congestion within the queue of the intermediate processor that trickles back to the windowed queuing unit 702 of the link layer processor 630. As such, downstream congestion trickles back upstream until windowed queue structure is reached, and at that point the windowing is used to ensure that this congestion is not increased.

Once the resource and routing processor 610 completes processing of data from its ingress queue 722, acknowledgements of such actions are typically generated that are to be sent back out in an egress message made up of one or more egress data units. These egress data units are forwarded to an egress queue in one of the plurality of intermediate processors, such as the egress queue 714 of the intermediate processor 622. Such acknowledgement messages are given preferential treatment by the intermediate processor 622 such that they are processed quickly. If the bandwidth within the intermediate processor 622 is limited, and the queue 714 begins to become filled, more bandwidth will be allocated to processing the acknowledgment messages than that allocated to processing the ingress call setup requests that may be pending in the ingress queue 712. This may result in the data in the ingress queue 712 exceeding the threshold level such that the windowing available at the queue 702 in the link layer processor 630 is used to ensure further congestion does not occur. As such, the high priority given to

processing of acknowledgment messages stored in the queue 714 ensures that congestion does not occur on the egress path through the various intermediate processors.

A similar strategy is used with respect to the egress queue 706 included within the link layer processors 630. By giving egress data units processing priority within the link layer processor 630, the windowing function available on the queue 702 can help to
5 ensure that congestion does not overwhelm the link layer processor 630.

The link layer processor 630 may also include a transmit queue 704 that is used to transmit acknowledgement messages and other egress data units. The transmit queue may utilize a sliding window that allows for multiple messages to be passed without
10 acknowledgement. When acknowledgment for any of those messages is received, another message can be sent such that there are always a certain number of messages that have been sent and are awaiting acknowledgement. For example, if the sliding window allows for a window size of three messages, and five messages have been received for transmission, three may be initially sent. The other two that have not yet been sent are
15 stored in the transmit queue until an acknowledgement is received corresponding to one of the initial three that were sent out. Once acknowledgement is received for one of the original three messages sent out, then one of the remaining two stored within the transmit queue can be sent out.

The transmit queue 704 of the link layer processor 630 transmits egress data units
20 to a downstream node, where the downstream node receives these egress data units as ingress data units at a queue similar to the queue 702. Thus, the downstream node will have a windowed queue structure, where the windowed queue may limit the transmission bandwidth available to the transmit queue 704. As such, if the downstream node becomes congested, the transmit queue 704 may become full. In such cases, some
25 messages stored within the transmit queue for transmission may have to be discarded in order to ensure that other more important messages are not lost. In some embodiments, the decision to discard egress data units may be made by one or more of the intermediate

processors such that discarded egress data units never reach the link layer processor 630. In other embodiments, the decision to discard egress data units may be made by the link layer processor 630 such that the data units are discarded after receipt from the intermediate processors.

5 Preferably, such message discarding by the transmit queue 704 is performed in an intelligent manner such that the consequences of such discard actions are minimized. For example, if new call setup messages are being forwarded through the network and are present within the transmit queue 704 along with acknowledgement messages corresponding to calls that have been recently established, the new call setup messages
10 are preferentially discarded before any acknowledgement messages are discarded. A good deal of processing bandwidth may have already been expended to generate the conditions resulting in the acknowledgement message such that discarding such an acknowledged message would effectively waste all of those resources that have already been expended. In the case of a call setup message, the resources may not yet have been
15 expended, and a subsequent call setup message can be issued without causing the high level of waste associated with discarding an acknowledgment message.

Once the congestion in the transmit queue 704 builds up to a threshold level, notification can be provided to the other processors within the switch such that those processors can make intelligent decisions regarding routing or other functions that may
20 help to alleviate the congested condition at the transmit queue 704. This may include rerouting of calls or rejection of call attempts by the resource routing processor. Note that each port in a line card may have a similar transmit queue associated with it such that there are a number of transmit queues within the switch. Therefore, those that have become congested may be avoided by these intelligent routing decisions such that the
25 congested transmit queues are allowed to recover from their congested condition.

Additional data paths exist with respect to the messaging processor 640 and the line card 652. Messages directed from the intermediate processor 622 to the ingress

queue 732 of the message processor 640 may correspond to programming commands that are directed towards the various line cards that are supported by the message processor 640, including the line card 652. As the message processor 640 receives such ingress messages at the queuing point 732, threshold detection is used to ensure that the queuing
5 point 732 does not become overly congested. If that threshold is exceeded, an indication is provided to the intermediate processor 622 and the resource and routing processor 610 so intelligent decisions can be made by routing and queue services such that generation of additional messages being fed to the message processor 640 is reduced or halted. This allows the congestion at the queue 742 to be relieved, and may result in the trickling
10 down of such congestion eventually to the windowed queue 702.

In most cases, congestion at the queuing point 742 will not occur due to the fact that the messaging processor is typically a reasonably fast processor that performs fairly simplistic operations. As such, it is typically capable of providing the level of bandwidth needed to keep up with the distribution of messages to the various line cards and
15 responses to the intermediate processors.

Within the message processor 640 there is a transmit queue associated with each of the line cards supported by the message processor 640. Thus, in the example shown in Figure 7, the transmit queue 734 supports the line card 652. The transmit queue 734 includes a threshold value that, when exceeded, acts in a similar manner as the transmit
20 queue 704 of the link layer processor 630. Thus, when the threshold is exceeded, the other processors included within the switch are notified. This may allow the processors to perform intelligent reroutes, stop changing the configuration of specific line cards, or do other things that help alleviate the congested condition. Preferably, no discarding of messages directed towards the line cards occurs within the transmit queue 734. This is
25 because the messages directed to the line cards are associated with programming or deprogramming operations that result in a changed configuration of different devices included in the line cards. If such messages are discarded, the state of the line cards may not be properly established, thus resulting in errors within the network.

The queue 742 in the line card 652 is preferably a windowed queue such that it can control the inflow of data from the transmit queue 724 of the message processor 640. Once a configuration message is processed by the line card 652, an acknowledgement message is typically returned back to the intermediate processor 622. Because the inflow of configuration messages to the line card 652 is limited by the queue 742, the outflow of acknowledgement messages is similarly limited. However, the queue 736 of the message processor 640 may include threshold detection such that if the threshold is exceeded, the window corresponding to queue 742 may be closed until the acknowledgements sent to queue 736 are dealt with properly. Acknowledgements relayed from the queue 736 of the message processor 640 to the queue 716 of the intermediate processor 622 are given high priority within the intermediate processor 622. Thus, if congestion begins to occur within the intermediate processor 622, additional bandwidth is provided to processing the acknowledgements such that bandwidth is taken away from processing ingress data units in the queue 712. As such, the congestion may trickle down to the windowed queue 702 for eventual alleviation.

The various thresholds associated with the queues included in the routing system of the communication switch are preferably configured to a reasonably low level such that a steady state condition will generally exist along the data paths of the multiprocessing system block. Higher thresholds for these queues might allow for concentrations in traffic that could overwhelm a particular queue when that queue is too far away from a more controlled queuing point.

In other embodiments, the transmit queues for each of the line cards within the message processor may be moved into various intermediate processors in the plurality of intermediate processors included within the switch. As such, the feedback path corresponding to the provision of messages from the intermediate processor 622 to the line card 652 would exist within the intermediate processor 622. If congestion were detected, a broadcast of this congestion to the other processors within the system is still possible, however, centralizing such transmit queue structures within the message

processor 640 provides added simplicity with respect to managing these transmit queues associated with the various line cards included within the switch.

By being able to detect a backlog of configuration messages directed toward the line cards in the transmit queues of the message processor 640, information regarding such congestion can be relayed to the various processors included within the switch. This allows the processors to make intelligent routing decisions that can allow such congestion to be alleviated, and also ensure that further congestion within these transmit queues does not occur such that messages aren't delivered within acceptable time periods. These features are not present in prior art systems, and become increasingly useful as data communication data rates continue to increase.

By coupling the ability to detect congestion at certain portions of the switch within a centralized routing processor, a much higher probability for determination of alternate routes based on congested conditions exists. In prior art systems, even if congested situations are detected, the potential for finding an appropriate alternate route is not as high as the routing and resource allocation processing within such prior art systems is distributed amongst the various processors included in each of the individual line cards. As such, these individual processors have a limited perspective and are only presented with a fixed set of potential ultimate routes if they were in fact notified of such congestion on a particular output port. Within the system described herein, the number of alternate routes is greatly increased due to the global viewpoint of the centralized resource allocation and routing.

In the foregoing specification, the invention has been described with reference to specific embodiments. However, one of ordinary skill in the art appreciates that various modifications and changes can be made without departing from the scope of the present invention as set forth in the claims below. Accordingly, the specification and figures are to be regarded in an illustrative rather than a restrictive sense, and all such modifications are intended to be included within the scope of present invention.

Benefits, other advantages, and solutions to problems have been described above with regard to specific embodiments. However, the benefits, advantages, solutions to problems, and any element(s) that may cause any benefit, advantage, or solution to occur or become more pronounced are not to be construed as a critical, required, or essential
5 feature or element of any or all the claims. As used herein, the terms "comprises," "comprising," or any other variation thereof, are intended to cover a non-exclusive inclusion, such that a process, method, article, or apparatus that comprises a list of elements does not include only those elements but may include other elements not expressly listed or inherent to such process, method, article, or apparatus.

CLAIMS

What is claimed is:

1. A buffering system in a communication switch, comprising:

5 a multiprocessor control block that includes:

a plurality of distributed processors that include ingress and egress
queuing points corresponding to data units communicated within the
communication switch, wherein when a congestion condition exists at selected
queuing points within the a distributed processor, a congestion indication is
10 generated; and

a resource routing processor operably coupled to the plurality of
distributed processors, wherein the resource routing processor controls routing
functionality within the communication switch, wherein the resource routing
processor receives congestion indications and preferentially selects uncongested
15 routes for subsequent connections within the communication switch based on the
congestion indications.

2. The buffering system of claim 1, wherein the resource routing processor performs
resource allocation amongst connections supported by the switch.

3. The buffering system of claim 2, wherein the plurality of distributed processors
includes:

a plurality of intermediate processors operably coupled to the resource routing
processor, wherein each intermediate processor of the plurality of intermediate processors
25 performs call processing for a corresponding portion of the connections supported by the

switch, wherein call processing includes issuing resource allocation requests to the resource routing processor, wherein each intermediate processor of the plurality of intermediate processors performs functions associated with a signaling layer portion of a protocol stack; and

5 a link layer processor operably coupled to the plurality of intermediate processors, wherein the link layer processor is operable to couple to a switching fabric of the communication switch, wherein the link layer processor receives ingress data units from the switching fabric and selectively forwards each ingress data unit received to at least one of the plurality of intermediate processors, wherein the link layer processor receives
10 egress data units from the plurality of intermediate processors and forwards each of the egress data units to the switching fabric.

4. The buffering system of claim 3, wherein the link layer processor includes a
15 windowing function such that the link layer processor controls the rate of receipt of ingress data units.

5. The buffering system of claim 4, wherein each intermediate processor of the plurality of intermediate processors includes an ingress buffer that buffers ingress data units received from the link layer processor, wherein when a threshold level of the
20 ingress buffer is exceeded, a threshold violation indication is generated that is provided to the link layer processor such that the link layer processor reduces flow of ingress data units to the ingress buffer whose threshold has been exceeded.

6. The buffering system of claim 5, wherein each intermediate processor of the
25 plurality of intermediate processors receives egress data units from the resource routing processor, wherein each intermediate processor of the plurality of intermediate processors preferentially processes egress data units with respect to ingress data units such that congestion in the intermediate processor is isolated to the ingress buffer.

7. The buffering system of claim 6, wherein the link layer processor receives egress data units from the plurality of intermediate processors, wherein the link layer preferentially processes egress data units with respect to ingress data units.

5 8. The buffering system of claim 7, wherein the link layer processor includes a transmit queue that buffers egress data units prior to transmission, wherein the transmit queue is a selected queuing point such that when the transmit queue becomes congested, the link layer processor generates a congestion indication that is provided to the resource routing processor.

10 9. The buffering system of claim 8, wherein when capacity of the transmit queue is exceeded, the link layer processor selectively discards egress data units based on a predetermined discard scheme.

15 10. The buffering system of claim 8, wherein when capacity of the transmit queue is exceeded, at least one intermediate processor of the plurality of intermediate processors selectively discards egress data units based on a predetermined discard scheme.

20 11. The buffering system of claim 8, wherein when capacity of the transmit queue is exceeded, the resource routing processor performs at least one of: rejecting a call attempt and selecting an alternate route for a call.

12. The buffering system of claim 2 further comprises:

25 a plurality of line cards operably coupled to the multiprocessor control block, wherein the plurality of line cards include ingress and egress queuing points for line card data units, wherein when a congestion condition exists at a queuing point within a line card, a line card congestion indication is generated and provided to the resource routing processor such that the resource routing processor selects routes at least partially based
30 on line card congestion indications received.

13. The buffering system of claim 12 further comprises:

a message processor operably coupled to the multiprocessor control block and the plurality of line cards, wherein the message processor supports messaging between the plurality of intermediate processors and the plurality of line cards.

14. The buffering system of claim 13, wherein the message processor includes an egress buffer that buffers egress line card data units received from the plurality of intermediate processors, wherein when a threshold level of the egress buffer is exceeded, an messaging threshold violation is generated.

15. The buffering system of claim 14, wherein the message processor includes a plurality of line card transmission queues, wherein each line card transmission queue of the plurality of line card transmission queues corresponds to one line card of the plurality of line card transmission queues, wherein each line card transmission queue buffers egress line cards data units directed to a corresponding line card, wherein when a line card transmission queue become congested, a line card congestion indication is generated and provided to the resource routing processor.

16. The buffering system of claim 15, wherein each line card of the plurality of egress line cards includes a windowing function such that the line card controls the rate of receipt of egress line card data units from a corresponding line card transmission queue in the message processor.

17. A communication switch, comprising:

a routing control block that performs call processing operations within the communication switch;

5

a plurality of line cards operably couple to the routing control block, wherein each of the line cards includes at least one transmit queue, wherein when congestion is detected on a transmit queue, a congestion indication is provided to the routing control block such that calls are routed away from the congestion.

10

18. The communication switch of claim 17, wherein the routing control block includes a plurality of processors, wherein each processor of the plurality of processors is responsible for a portion of the protocol stack used in call processing operations, wherein each processor includes queuing points, wherein a first set of queuing points of the queuing points in the communication switch are rate controlled in a manner that ensures that congestion at the first set of queuing points does not occur, wherein when congestion occurs and is detected at queuing points included in a second set of queuing points, notification is provided to a routing processor of the plurality of processors, wherein the routing processor performs subsequent routing operations based on congestion notifications.

15
20

19. A method for call processing in a communication switch, comprising:

detecting congestion in a transmit queue corresponding to a line card of the communication switch; and

5

providing an indication of the congestion to a central control block that performs call processing and routing for a plurality of line cards included in the communication switch, wherein the central control block performs subsequent routing operations in a manner that avoids the congestion corresponding to the line card.

10

20. The method of claim 19, wherein the central control block includes a resource routing processor, a plurality of intermediate processors, and a link layer processor, wherein the resource routing processor performs the subsequent routing operations.

15

21. The method of claim 19, wherein performing subsequent routing operations includes maintaining status of a plurality of transmit queues corresponding to a plurality of line cards in the switch, wherein the status is used to determine a non-congested compatible transmit queues for the subsequent routing operations.

20

22. The method of claim 21 further comprises prioritizing data flow in the switch such that congestion is concentrated at the plurality of transmit queues.

25

23. The method of claim 19, wherein the congestion in the transmit queue is a result of a buildup of messages corresponding to programming commands that are directed towards the line card.

**BUFFERING SYSTEM FOR USE IN A COMMUNICATION SWITCH THAT
INCLUDES A MULTIPROCESSOR CONTROL BLOCK AND METHOD
THEREFORE**

5

Abstract of the Invention

A buffering system in a communication switch and method therefore is presented that utilizes notification of congested points within the switch to perform intelligent
10 routing decisions such that the congestion is avoided when possible. In one embodiment, this is accomplished by detecting congestion in a transmit queue corresponding to a line card or a link layer processor of the communication switch. Upon detection of such congestion, an indication of the congestion is provided to a central control block that performs call processing and routing for the communication switch. When the central
15 control block performs subsequent routing operations, the central control block considers the congestion indications it has thus far received such that the congested points within the switch are avoided if possible. Techniques are also presented for isolating congestion within the switch to predetermined congestion points, where when congestion occurs at these predetermined points, specific circuitry is included to help ensure that the
20 congestion is controlled.

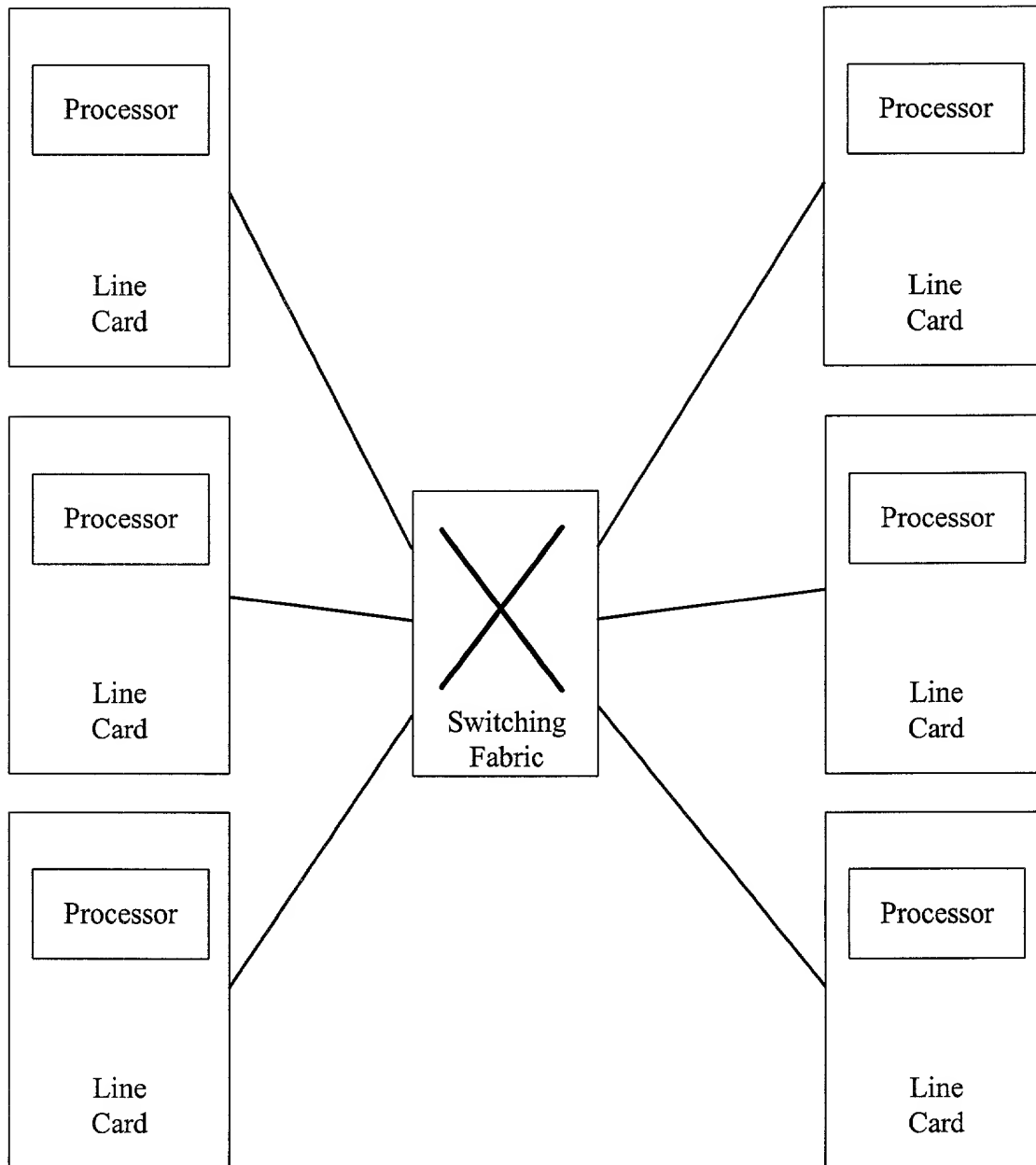


Figure 1.
(PRIOR ART)

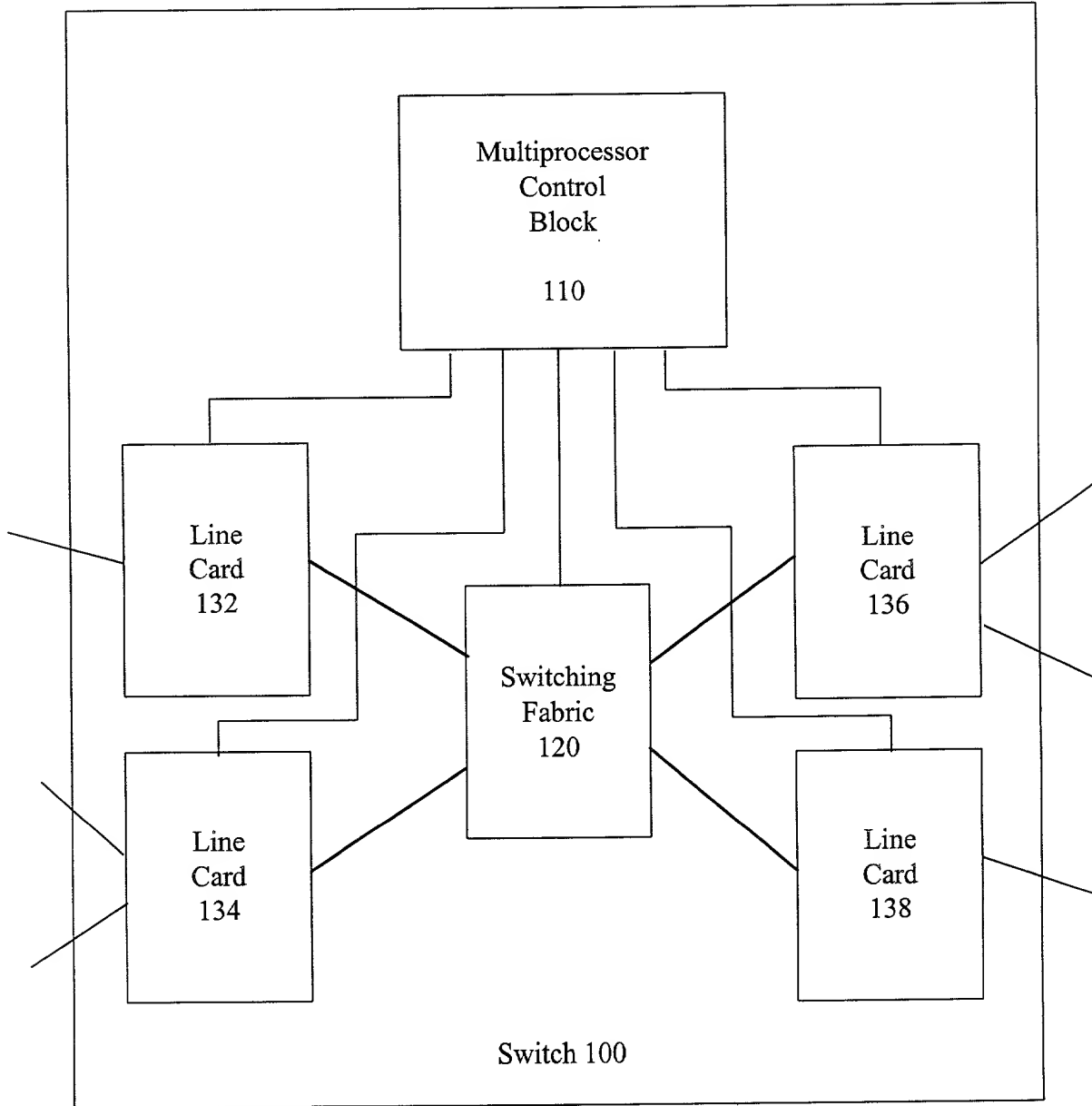


Figure 2.

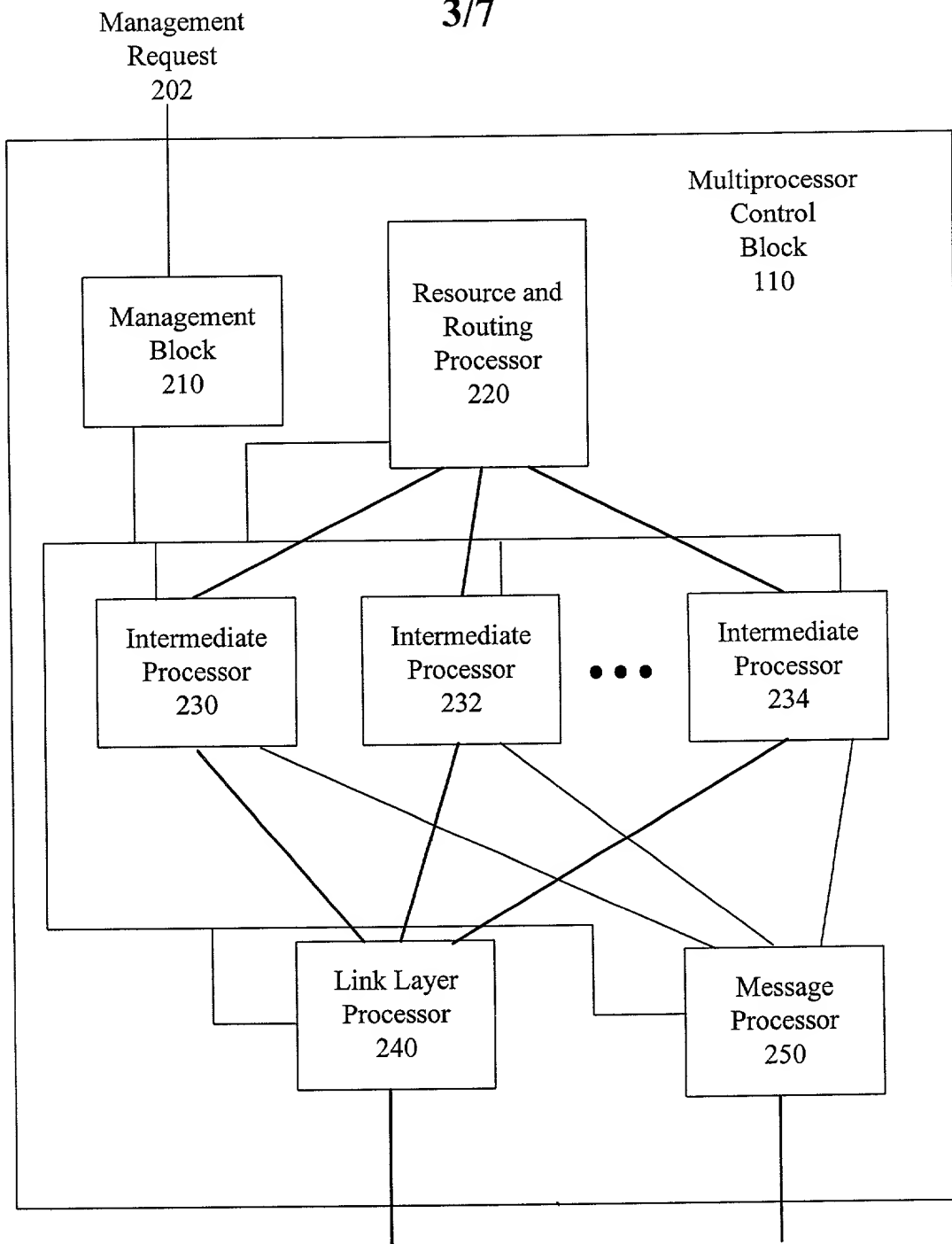


Figure 3.

Protocol Stack 290

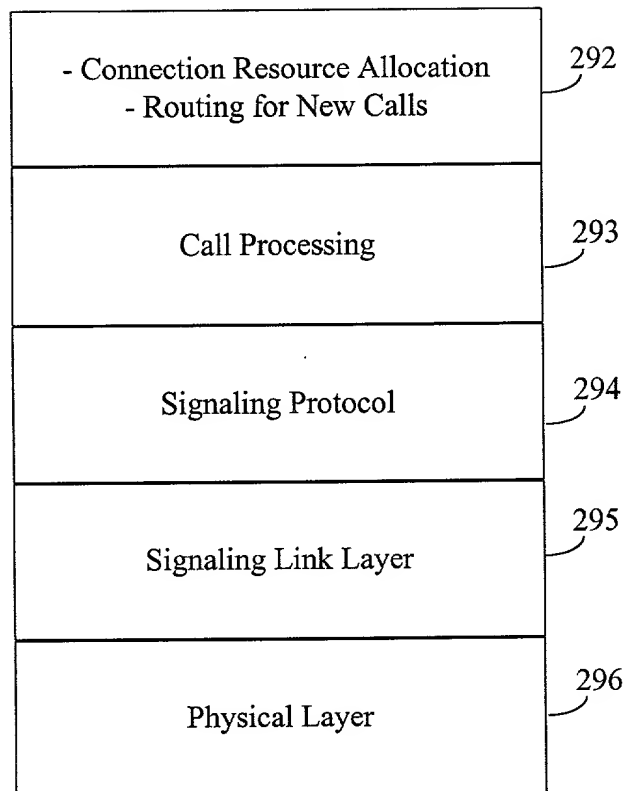


Figure 4.

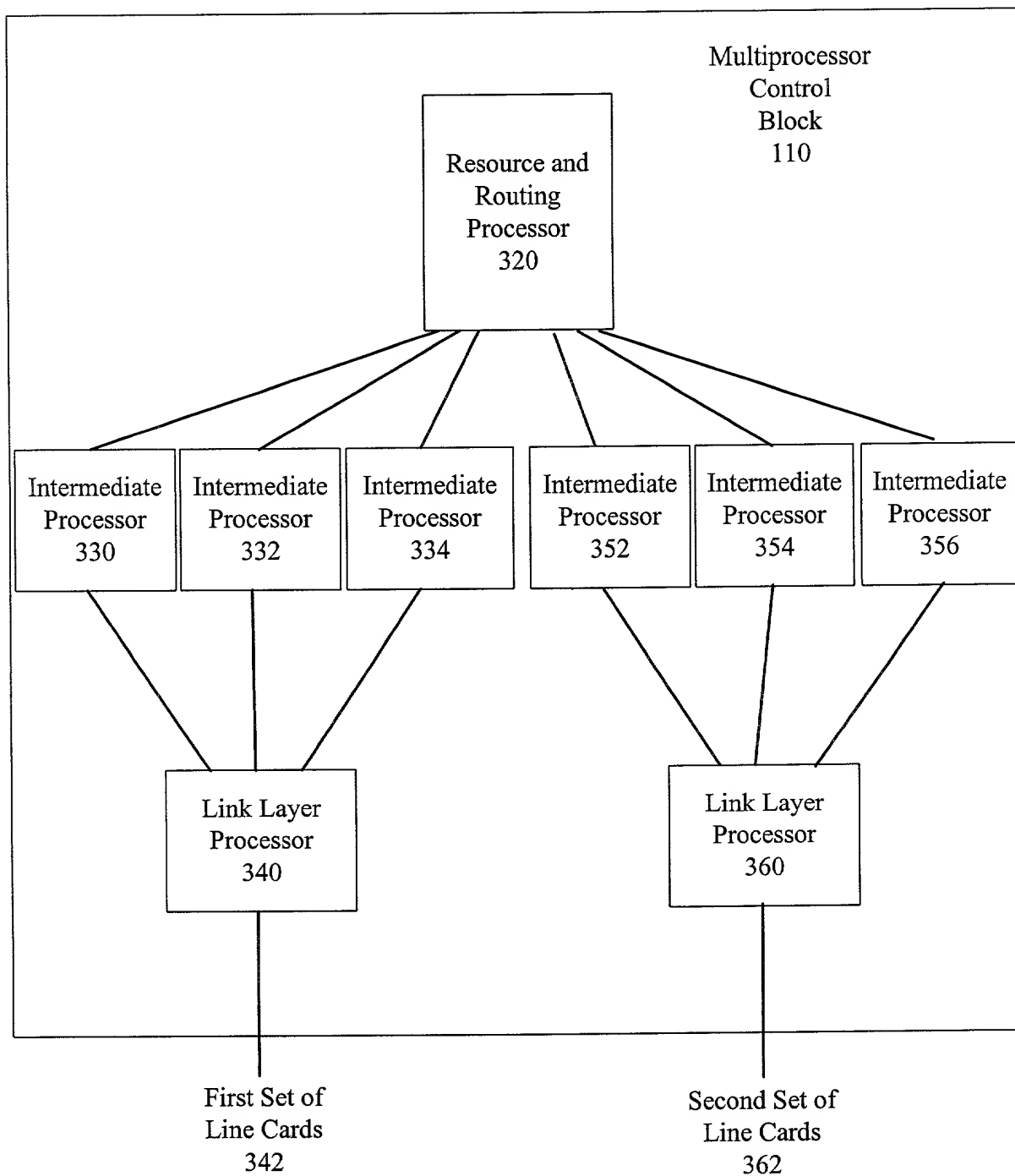


Figure 5.

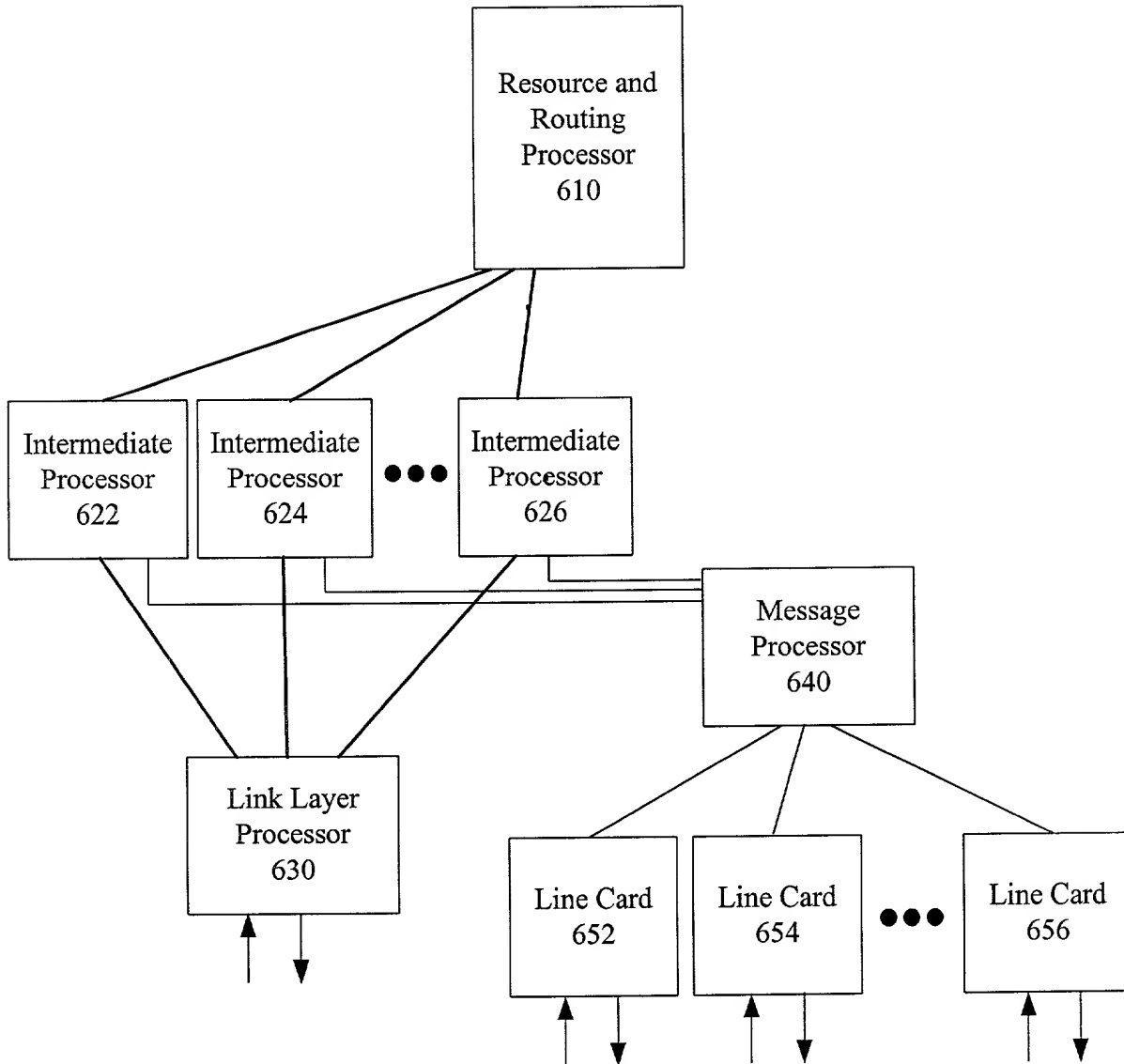


Figure 6.

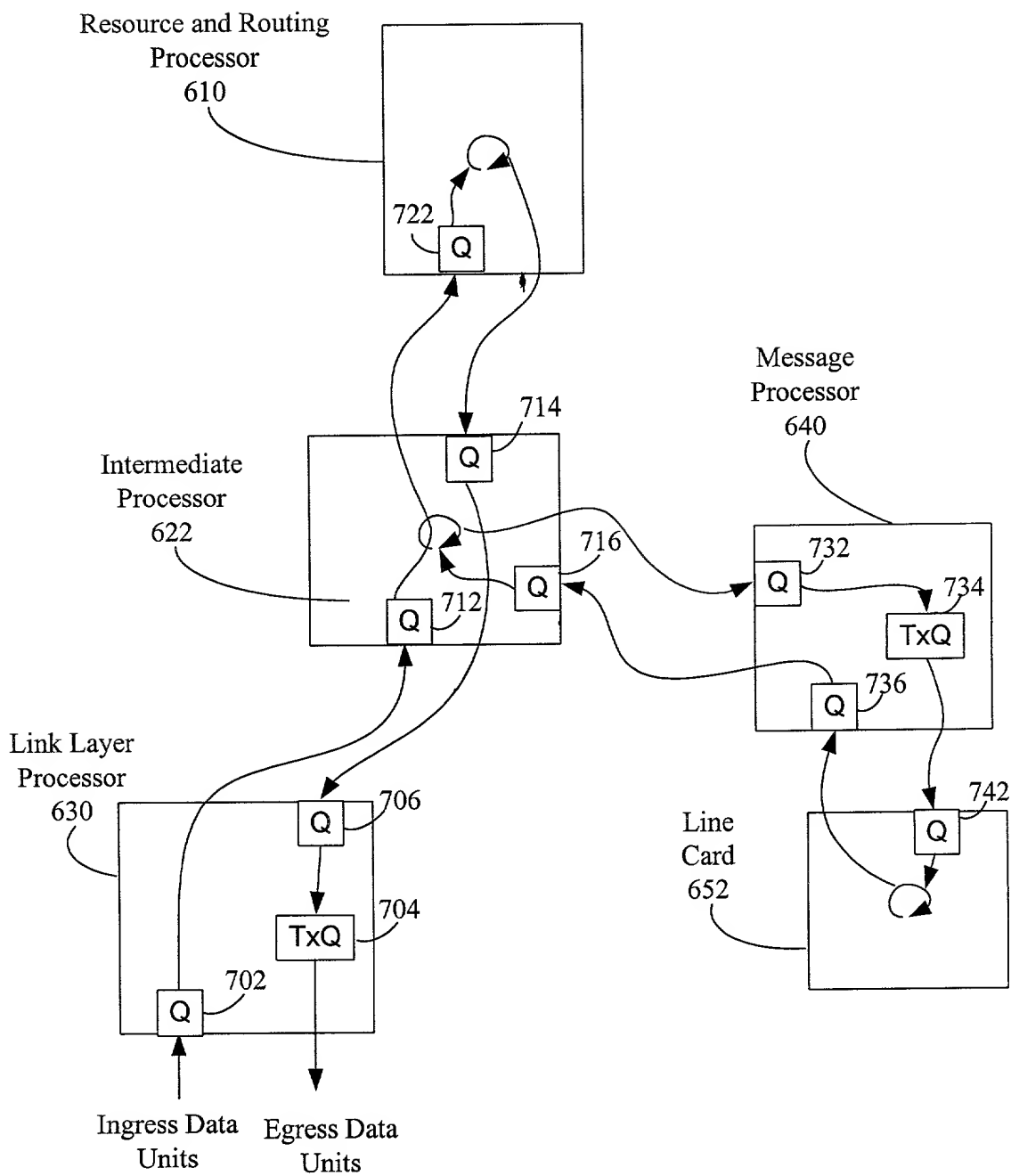


Figure 7.

**DECLARATION
FOR UTILITY OR DESIGN
PATENT APPLICATION
(37 CFR 1.63)**

- ☒ Declaration Submitted with Initial Filing, OR
☐ Declaration Submitted after Initial Filing
(surcharge (37 CFR 1.16 (e)) required)

Attorney Docket Number 1400.4100290
First Named Inventor McCormick, et al.
COMPLETE IF KNOWN
Application Number:
Filing Date:
Group Art Unit :
Examiner Name

As a below named inventor, I hereby declare that:

My residence, post office address, and citizenship are as stated below next to my name.

I believe I am the original, first and sole inventor (if only one name is listed below) or an original, first and joint inventor (if plural names are listed below) of the subject matter which is claimed and for which a patent is sought on the invention entitled:

**BUFFERING SYSTEM FOR USE IN A COMMUNICATION SWITCH THAT INCLUDES
A MULTIPROCESSOR CONTROL BLOCK AND METHOD THEREFORE**

the specification of which:

- ☒ is attached hereto.
☐ was filed on (MM/DD/YYYY) as United States Application Number or PCT International Application Number and was amended on (MM/DD/YYYY) (if applicable).

I hereby state that I have reviewed and understand the contents of the above identified specification, including the claims, as amended by any amendment specifically referred to above.

I acknowledge the duty to disclose information which is material to patentability as defined in 37 CFR 1.56.

I hereby claim foreign priority benefits under 35 U.S.C. 119(a)-(d) or 365(b) of any foreign application(s) for patent or inventor's certificate, or 365(a) of any PCT international application which designated at least one country other than the United States of America, listed below and have also identified below, by checking the box, any foreign application for patent or inventor's certificate, or of any PCT international application having a filing date before that of the application on which priority is claimed.

Prior Foreign Application Number(s)	Country	Foreign Filing Date (MM/DD/YYYY)	Priority Not Claimed	Certified Copy Attached?	
				YES	NO
			<input type="checkbox"/>	<input type="checkbox"/>	<input type="checkbox"/>
			<input type="checkbox"/>	<input type="checkbox"/>	<input type="checkbox"/>

- ☐ Additional foreign application numbers are listed on a supplemental priority data sheet PTO/SB/02B attached hereto.

I hereby claim the benefit under 35 U.S.C. 119(e) of any United States provisional application(s) listed below.

Application Number(s)	Filing Date (MM/DD/YYYY)
60/224,441	08-10-2000

- ☐ Additional provisional application numbers are listed on a supplemental priority data sheet PTO/SB/02B attached hereto.

I hereby claim the benefit under 35 U.S.C. 120 of any United States application(s), or 365(c) of any PCT international application designating the United States of America, listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in the prior United States or PCT International application in the manner provided by the first paragraph of 35 U.S.C. 112, I acknowledge the duty to disclose information which is material to patentability as defined in 37 CFR 1.56 which became available between the filing date of the prior application and the national or PCT international filing date of this application.

U.S. Parent Application or PCT Parent Number	Parent Filing Date (MM/DD/YYYY)	Parent Patent Number (if applicable)

- ☐ Additional U.S. or PCT international application numbers are listed on a supplemental priority data sheet PTO/SB/02B attached hereto.

As a named inventor, I hereby appoint the following registered practitioner(s) to prosecute this application and to transact all business in the Patent and Trademark Office connected therewith:

Name	Registration Number	Name	Registration Number
Christopher J. Reckamp	34,414		
Paul M. Anderson	39,896		
Ross D. Snyder	37,730		
John R. Garrett	27,888		

☐ Additional registered practitioner(s) named on supplemental Registered Practitioner Information sheet PTO/SB/02C attached hereto.

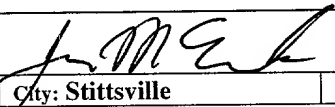
Direct all correspondence to:

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115 Wild Basin Road-Suite 107
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Facsimile: 512-347-9224

I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under 18 U.S.C. 1001 and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.

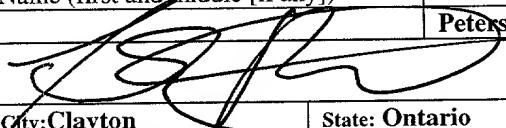
Name of Sole or First Inventor:

☐ A petition has been filed for this unsigned inventor

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Name of Additional Joint Inventor:

☐ A petition has been filed for this unsigned inventor

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☐ Additional inventors are being named on the _____ supplemental Additional Inventor(s) sheet(s) PTO/SB/02A attached hereto.